Math 564: Real analysis and measure theory Lecture 8

Terminology. Let (X, B, M) be a measure space and P be a property of points in X (e.g. being transcendental for $\chi := \mathbb{R}$). Then we say that

Pholos a.e. (almost everywhere) in X) or a.e. (almorst every) x & X satisfies P y it \x & X: x satisfies P > is conull. or Pholds a.s. (almost shrely)

Measure exhaustion.

In a measure space, call a collection C of sets almost disjoint if the pairwise taker-uctions of sets in C are well.

almost disjoint collection & of μ -measures). Let (x, B, μ) be a refinite measure space. Then any

Proof. We first prove this assuming $\mu(X) < \infty$. Then for each $\mu \in (N^4)$, the self C, = { C ∈ C : μ(c) > { }

is timbe (in tact, zn.p(x) elements) and Z= V, so E is of timbe wessive.
For the general T= Finite case, let X= || Xn where cach Xn \in B is of finite messive.
Acil define

 $\mathcal{D}_{n} := \left\{ C \in \mathcal{C} : \mu(C \cap X_{n}) > 0 \right\}.$ Then by the tinide case, each \mathcal{D}_{n} is often and $\mathcal{C} = \mathcal{U} \mathcal{D}_{n}$, so \mathcal{C} is often.

Transfinite measure exhaustion. Let (X,B,μ) be a σ -finite measure space and let $(A_{\alpha})_{\alpha}<\omega$, be an increasing represent of μ -measurable sets, where ω , is the first most ordinal. Then the segmence almost stabilizes at some of the ordinal Υ , i.e. $\forall d \gg \Upsilon$, $A_{\alpha}=\mu A_{\alpha}$.

Remark This allows to sun transfirst algorithms which at each step handle a positive measure set. Then we know he algorithm will stop at a cottol stage, having handled a countly set.

We now discuss an important application. In a measure space with atoms, we can't achieve every value of measure between O and p(X), but this is the only obstraction.

Sierpinski's theorem. In an atomless measure space (K, B, N), every value 0 < r < p(X) is a diseved, i.e. here is B = B with p(B) = r.

Proof First let's prove a more humble statement:

Claim I treey positive measure set Y contains positive measure sets of expitacily small measure.

Pf of Claim Y is not an atom so there must be $X \circ \in Y$ with $\mu(X_{0}) < \mu(Y)$. We bailed a sequence $\{X_{s}\}_{s \in 2^{c}N}$ of positive measure sets such that $X_{s} = X_{so} \sqcup X_{ss}$ as tollows: if X_{s} is defined, it's not an atom, so thore is $X_{so} \subseteq X_{s}$ in $X_{so} \sqcup X_{so}$ with $0 < \mu(X_{so}) < \mu(X_{s})$. Let $X_{s1} := X_{s} \setminus X_{so}$.

To each $s \in 2^{c}N$ one of X_{so} and X_{s1} has necessare $\leq \frac{1}{2}\mu(X_{s})$.

The which gives an infinite branch $(X_{sn})_{n \in N}$ in the tree of positive vectors $(X_{sn})_{n \in N}$ in the tree of positive vectors $(X_{sn})_{n \in N}$ in the tree of positive vectors.

Iteratively using Claim I, we now explicitly build a set B & B with p(B) = r.

Proof via transficite exhaustion. Delive a sequence (Aa) acros = B of pairwise disjoint sets such that $\mu(\Box A_a) \in \Gamma$ for each $\beta \in W_1$, by inclustion as follows: if $ A_a _{\alpha \in \beta}$ is already defined, let $ A_a _{\alpha \in \beta}$ be a positive recover subset of $X \setminus \bigcup_{\alpha \in \beta} A_{\alpha}$ of resource $\forall r - \mu(\Box A_{\alpha})$ if $r - \mu(\Box A_{\alpha}) = 0$ otherwise put $ A_{\alpha} = 0$. Now the proof of of the pigeonhole for measures (using the weelition $ A_{\alpha} = 0$) instead of the finiteness of $ A_{\alpha} $ gives that all but of $ A_{\alpha} = 0$.
Proof via $\frac{1}{2}$ -greedy algorithm. We inductively build a sequence $(B_n)_{n\in\mathbb{N}} \subseteq B$ of pairwise disjoint sets such that $\mu(\sqcup B_i) \leq r$ as follows: suppose $(B_i)_{i\geq n}$ is defined and take $B_n \in B$ to be any set with
Now that (Bu) new is defined, monotone convergence implies $\sum_{n \in \mathbb{N}} \mu(B_n) = \mu(A_n B_n) \leq r$, in particular, lim $\mu(B_n) = 0$. We now check that the set
has measure = Γ . Indeed, otherwise, $\mu(Bw) < r$, so by Claim 1, there is $B' \leq X \setminus Bw$ in B such that $0 < \mu(B') \leq r - \mu(Bw)$. But taking a large enough well so that $\mu(Bu) < \frac{1}{2} \mu(Bw)$
we get a voit adjution with the choice of Bn. Approximating measurable subs
- Into ki-wanta was a long street sold

99% lemma. We begin with a basic observation. Observation (percentage of carrots in soup). Let (X, B, p) be a measure space and let A, B be p-measure cable sets with $0 < \mu(B) < \infty$. Then for any (percentage) $p \in [0,1]$ and any (finite or clb1) partition $B = \bigcup_{n < N} B_n$, where $N \in N \cup \{\infty\}$, we have

$$\frac{\mu(A\cap B)}{\mu(B)} \ge p$$
 \Longrightarrow $\frac{\mu(A\cap B_n)}{\mu(B_n)} \ge p$ become nEIN.

 $\frac{\mu(A \cap B)}{\mu(B)} = \sum \frac{\mu(B_1)}{\mu(B)} \cdot \frac{\mu(A \cap B_n)}{\mu(B_n)}, \quad \text{when} \quad \sum \frac{\mu(B_n)}{\mu(B)} = 1, \text{ so it's a convex combination.}$